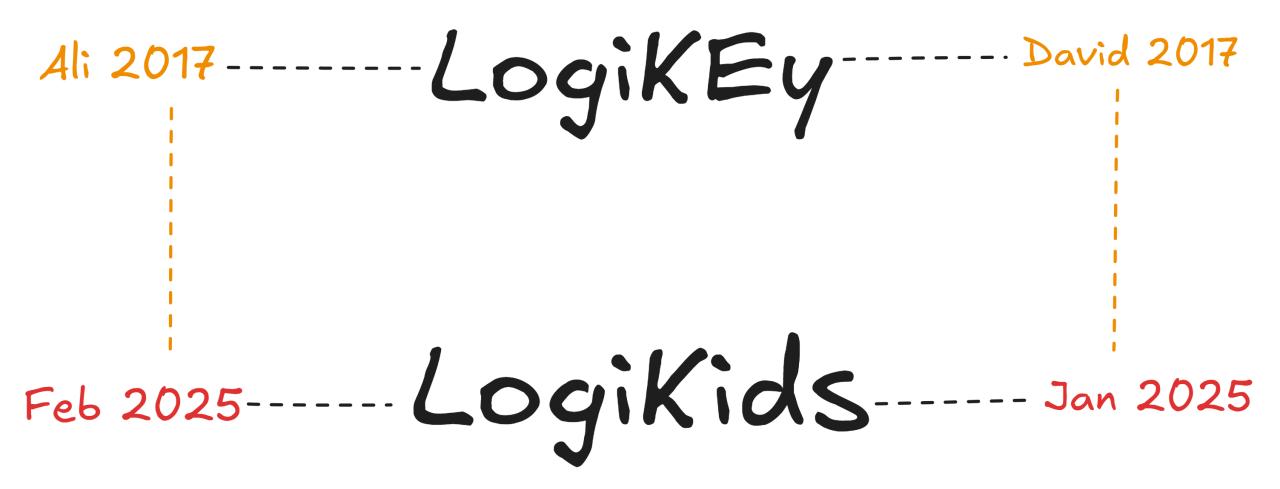
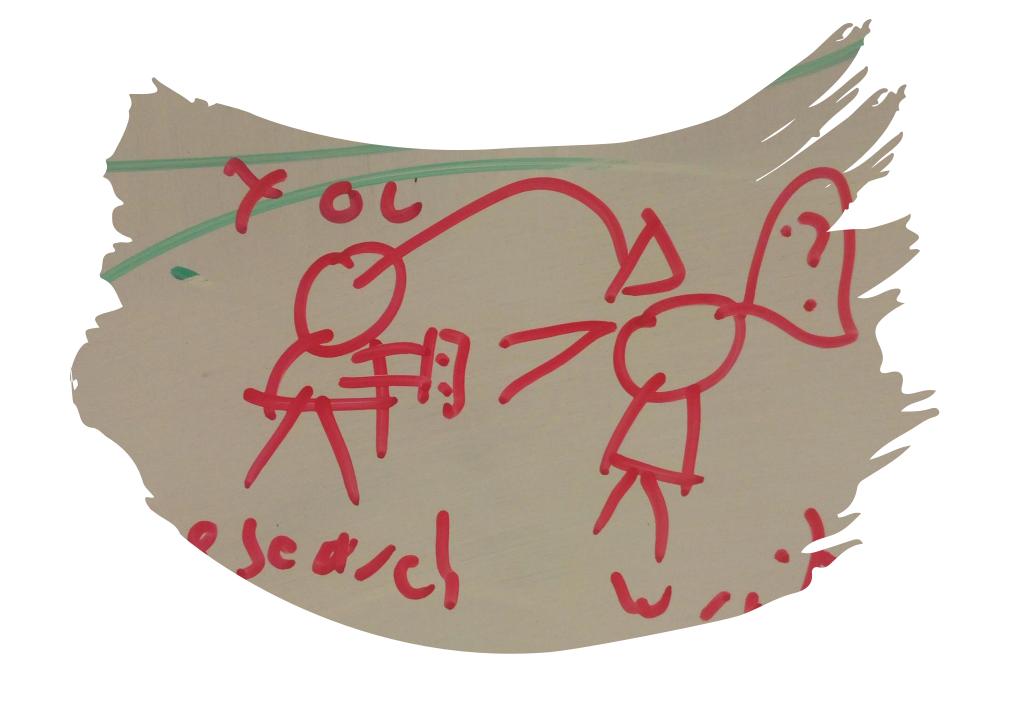
## LogikEy Plus: A Toolkit for Kids in 2030

Ali Farjami SnT



## 





By Stacey Frost ₳

Innovative AI partner for toy design.

Can you suggest a new toy design?

What are the latest trends in toy design?

Generate a unique toy idea.

design a random toy car.

## Mecha Brain Explorer

#### Obligations:

"You must light up when I press your buttons!"
"You must move when I push you!"

#### Permissions:

"You can make sounds when I play with you."
"You can go anywhere in my room."

#### Prohibitions:

"You must not break when I play with you."
"You cannot get dirty or lose your colours!"



The LogiKEy methodology is primarily a framework designed for ethical reasoning in complex systems, especially in AI, law, and governance. It focuses on using formal logic (particularly deontic logic) to model and solve ethical dilemmas by defining normative systems—rules or standards that govern behavior. This methodology is built on higher-order logic (HOL) and supports experimentation with various ethical theories, making it a flexible tool for reasoning about ethics in intelligent systems(FIS Uni Bamberg)(ar5iv).

While LogiKEy is mainly applied to areas like autonomous systems and legal reasoning, it could potentially be adapted for toy design if your goal is to ensure that your toy operates within clear ethical frameworks. For example, you might use LogiKEy to model the ethical implications of how a toy might influence behavior or decision-making in children. By applying LogiKEy, you could simulate scenarios where the toy encourages certain values (e.g., sharing, empathy) or analyze the fairness of its representations (inclusivity and diversity).

In practical terms, however, you might find that LogiKEy's approach is more useful for designing toys that include AI elements, such as interactive robots or smart educational toys. It would help ensure that these systems make ethically sound decisions when interacting with children or gathering data (ar5iv).

If you're not working with AI or autonomous systems in your toy design, other ethical design frameworks, like sustainable material usage or child development psychology, might be more directly applicable.

40 0 0 0

## LogiKEy

\*Obligation



\*Permission



\*Prohibition



\*Design Platform X

\*Online Update X

## LogikEy Plus

\*Obligation



\*Permission



\*Prohibition



\*Design Platform V



\*Online Update X



### LogiKEy -

\*Obligation



\*Permission



\*Prohibition



\*Design Platform V



\*Online Update



## LogiKEy

\*Obligation



\*Permission



\*Prohibition







#### Data in Brief

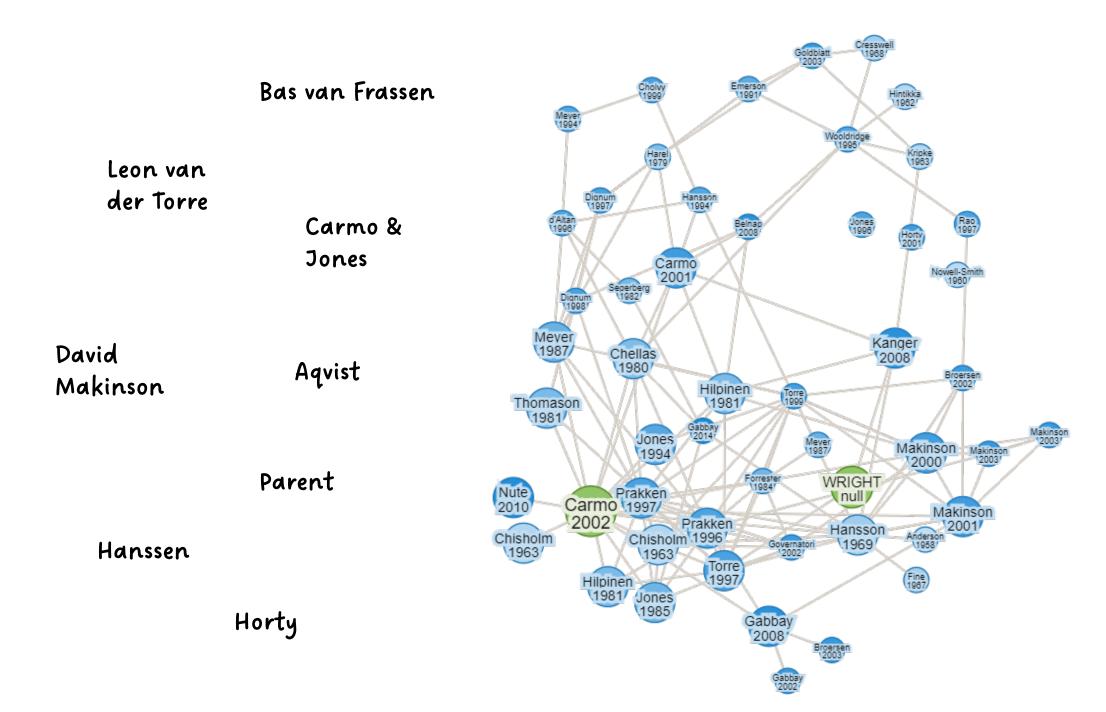
Volume 33, December 2020, 106409



Data Article

#### LogiKEy workbench: Deontic logics, logic combinations and expressive ethical and legal reasoning (Isabelle/HOL dataset)

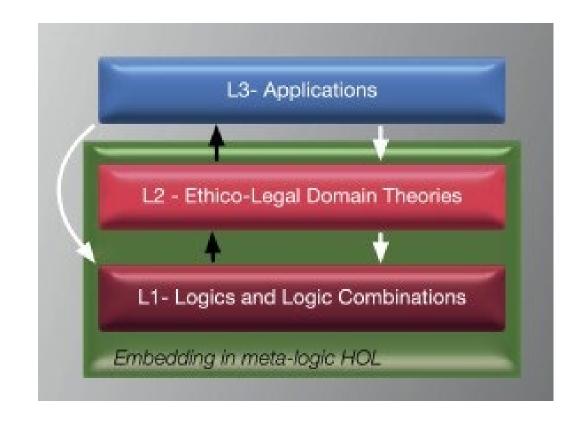
Christoph Benzmüller a b ス ⋈, Ali Farjami a, David Fuenmayor b, Paul Meder a, Xavier Parent <sup>a</sup> ス ⋈ , Alexander Steen <sup>a</sup>, Leendert van der Torre <sup>a c</sup>, Valeria Zahoransky <sup>a d</sup>

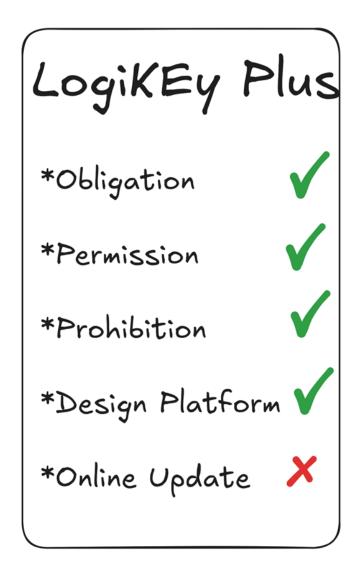


## We are focusing on

Layer L1

Making L1 Customizable





Thanks to:
Andrea De Domenico,
Krishna Manoorkar,
Alessandra Palmigiano,
Mattia Panettiere,
Xiaolong Wang

## Two upgrades:

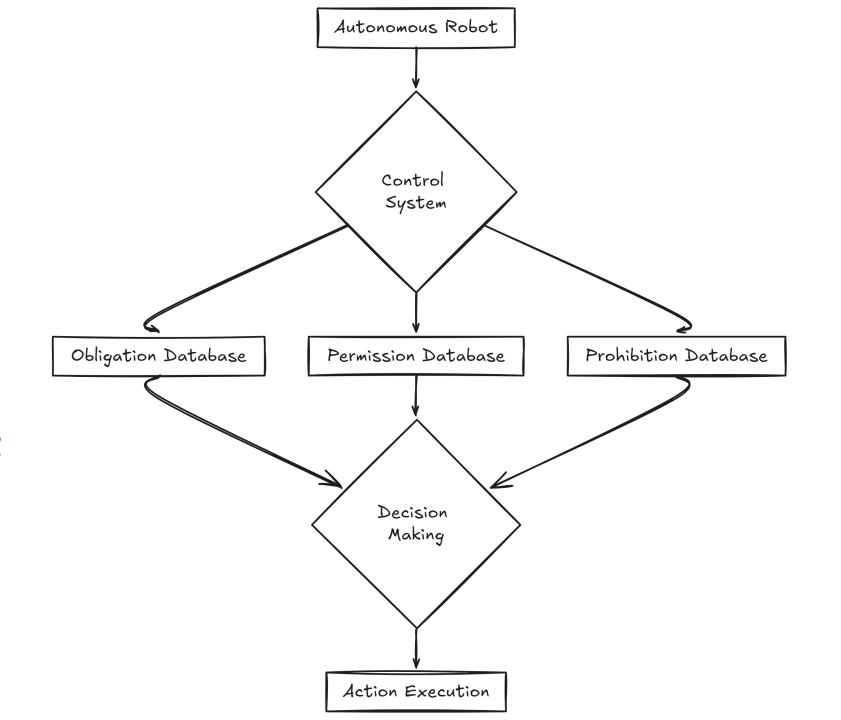
## Accessing Memory

Deontic Control Systems:

Conflict management Sorting and giving priority

#### Partial Database

No need for coordination between obligation, permission and prohibition



## Flexible Control System

## Input/Output Logic:

Free (Normative) Software Movement





Freedom of speech and expression Freedom of worship Freedom from want Freedom from fear

## Four Freedoms Speech



## Open Source:

## Easy to distribute:

The freedom to run the program as you wish

The freedom to study how the program works, and change it The freedom to redistribute copies

The freedom to distribute copies of your modified versions to others

### No Dark Pattern

Axiom schemata for propositional logic (PL) S5-schemata for 
$$\square$$
 and  $\diamondsuit$  (S5) 
$$\bigcirc (\psi_1 \to \psi_2/\varphi) \to (\bigcirc (\psi_1/\varphi) \to \bigcirc (\psi_2/\varphi))$$
 (COK) 
$$\bigcirc (\psi/\varphi) \to \square \bigcirc (\psi/\varphi)$$
 (Abs) 
$$\square \psi \to \bigcirc (\psi/\varphi)$$
 (Nec) 
$$\square (\varphi_1 \leftrightarrow \varphi_2) \to (\bigcirc (\psi/\varphi_1) \leftrightarrow \bigcirc (\psi/\varphi_2))$$
 (Ext) 
$$\bigcirc (\varphi/\varphi)$$
 (Id) 
$$\bigcirc (\psi/\varphi_1 \land \varphi_2) \to \bigcirc (\varphi_2 \to \psi/\varphi_1)$$
 (Sh) If  $\vdash \varphi$  and  $\vdash \varphi \to \psi$  then  $\vdash \psi$  (MP) If  $\vdash \varphi$  then  $\vdash \square \varphi$ 

## Four Freedoms for Deontic Logic

# You are free to choose your base logic:

Propositional, Intuitionistic, Modal,...

#### Obligations and permissions on selfextensional logics

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<sup>3</sup>Department of Mathematics and Applied Mathematics, University of Johannesburg, Corner Kingsway and University Road, Rossmore, Johannesburg, South Africa

School of Philosophy and Social Development, Shandong University, South Shanda Road No.27, Jinan, 250100, China

#### Abstract

We further develop the abstract algebraic logic approach to input/output logic initiated in [13], where the family of selfextensional logics was proposed as a general background environment for input/output logics. In this paper, we introduce and discuss the generalizations of several types of permission (negative, dual negative, static, dynamic), as well as their interactions with normative systems, to various families of selfextensional logics, thereby proposing a systematic approach to the definition of normative and permission systems on nonclassical propositional bases.

Keywords: input/output logic, selfextensional logics, abstract algebraic logic.

# You are free to fully characterize your logic:

Adding and Removing Axioms

#### Obligations and permissions, algebraically

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 <sup>c</sup>Department of Mathematics and Applied Mathematics, University of Johannesburg, South Africa,
 <sup>d</sup>School of Philosophy and Social Development, Shandong University, Jinan, China

#### Abstract

We further develop the algebraic approach to input/output logic initiated in [1], where subordination algebras and a family of their generalizations were proposed as a semantic environment of various input/output logics. In particular, we consider precontact algebras as a suitable algebraic environment for negative permission, and we characterize properties of several types of permission (negative, static, dynamic), as well as their interactions with normative systems, by means of suitable modal languages encoding outputs.

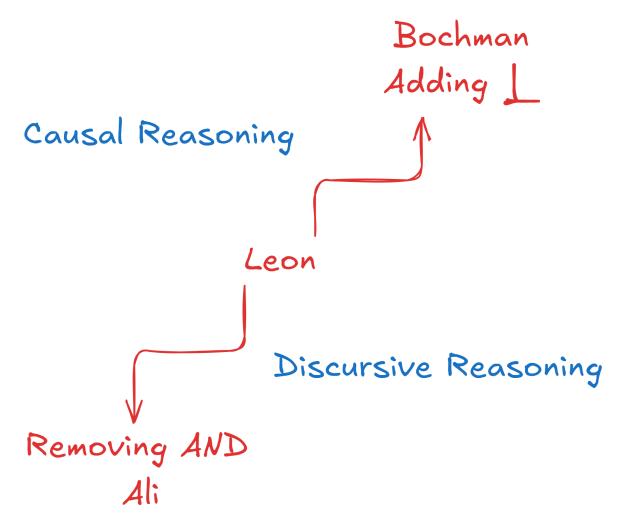
Keywords: input/output logic, subordination algebras, precontact algebras, selfextensional logics, slanted algebras, algorithmic correspondence theory.

2000 MSC: 03G25, 03G10

(L) 
$$\bot \prec \bot$$
 (T)  $\top \prec \top$   
(SB)  $\exists b(b \prec a)$  (SF)  $\exists b(a \prec b)$   
(SI)  $a \leq b \prec x \Rightarrow a \prec x$  (WO)  $b \prec x \leq y \Rightarrow b \prec y$   
(AND)  $a \prec x \& a \prec y \Rightarrow a \prec x \land y$  (OR)  $a \prec x \& b \prec x \Rightarrow a \lor b \prec x$   
(DD)  $a \prec x_1 \& a \prec x_2 \Rightarrow \exists x(a \prec x \& x \leq x_1 \& x \leq x_2)$ 

()	ce , te I	-	, 2		1 40				_ ~ 2	/
(UD)	$a_1 \prec x$	$\& a_2$	$\prec x \Rightarrow$	$\Rightarrow \exists a(a)$	$\prec x$	& a1	$\leq a$	& a	$2 \leq a$	

Name	Properties (DD) (SF)					
♦-defined						
-defined	(UD) (SB)					
defined	♦-defined + ■-defined					
◇-premonotone	$\diamond$ -defined + (SI)					
-premonotone	$\blacksquare$ -defined + (WO)					
◇-monotone	♦-premonotone + (WO)					
-monotone	$\blacksquare$ -premonotone + (SI)					
monotone	♦-monotone + ■-monotone					
♦-regular	$\diamond$ -monotone + (OR)					
<b>■</b> -regular	$\blacksquare$ -monotone + (AND)					
regular						
♦-normal	$\diamond$ -regular + ( $\perp$ )					
-normal	$\blacksquare$ -regular + $(\top)$					
subordination algebra	♦-normal + ■-normal					



# You are free to combine and design new logics

Kracht formulas: Ob + Per + Pro

#### Correspondence and Inverse Correspondence for Input/Output Logic and Region-Based Theories of Space

Andrea De Domenico<sup>a</sup>, Ali Farjami<sup>b</sup>, Krishna Manoorkar<sup>a</sup>, Alessandra Palmigiano<sup>a,c</sup>, Mattia Panettiere<sup>a</sup>, Xiaolong Wang\*<sup>a,d</sup>

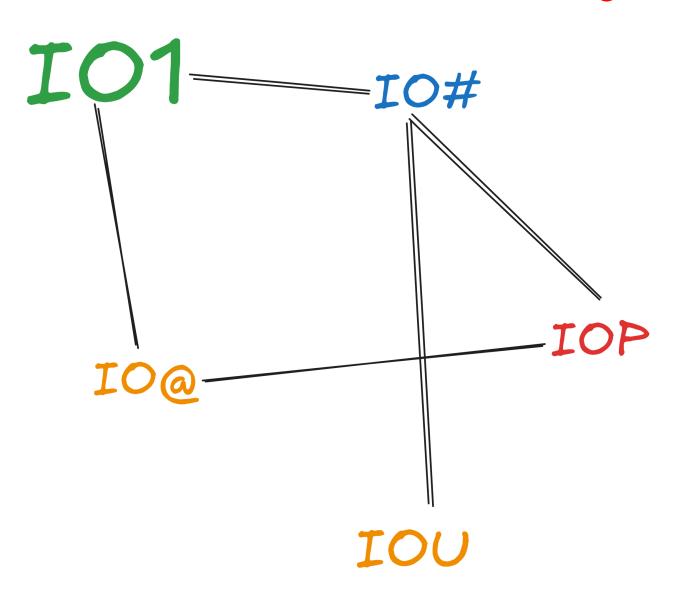
<sup>a</sup>School of Business and Economics, Vrije Universiteit Amsterdam, The Netherlands,
 <sup>b</sup>Interdisciplinary Centre for Security Reliability and Trust, University of Luxembourg, Luxembourg
 <sup>c</sup>Department of Mathematics and Applied Mathematics, University of Johannesburg, South Africa,
 <sup>d</sup>School of Philosophy and Social Development, Shandong University, Jinan, China

#### Abstract

We further develop the algebraic approach to input/output logic initiated in [27], where subordination algebras and a family of their generalizations were proposed as a semantic environment of various input/output logics. In particular: we extend the modal characterizations of a finite number of well known conditions on normative and permission systems, as well as on subordination, precontact, and dual precontact algebras developed in [29], to those corresponding to the infinite class of analytic inequalities in a modal language consisting both of positive and of negative unary modal operators; we characterize the syntactic shape of first-order conditions on algebras endowed with subordination, precontact, and dual precontact relations which guarantees these conditions to be the first-order correspondents of axioms in the modal language above; we introduce algorithms for computing the first-order correspondents of modal axioms on algebras endowed with subordination, precontact, and dual precontact relations, and conversely, for computing the modal axioms of which the conditions satisfying the suitable syntactic shape are the first-order correspondents; finally, we extend Celani's dual characterization results between subordination lattices and subordination spaces to a wider environment which also encompasses precontact and dual precontact relations, and relative to an infinite class of first order conditions relating subordination, precontact and dual precontact relations on distributive lattices.

Keywords: input/output logic, subordination algebras, subordination lattices, precontact algebras, Sahlqvist theory, inverse correspondence, Kracht's theorem.

## Lion Project



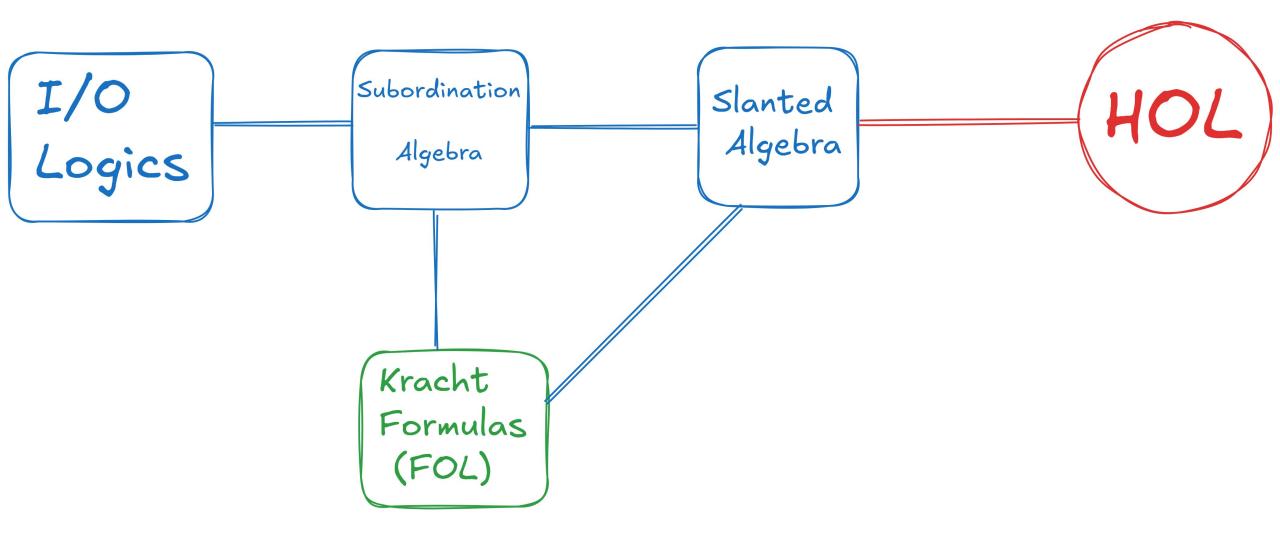
Dynamic Permission

Static Permission

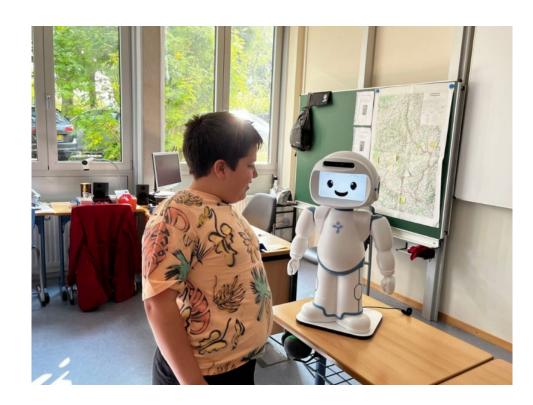
## You are free to efficiently implement your logic

```
consts IO2 :: "\tau \Rightarrow \tau" ("\diamondsuit^2<sub>0</sub>")
                       consts IO3 :: "\tau \Rightarrow \tau" ("\diamondsuit^3<sub>0</sub>")
                       consts I04 :: "\tau \Rightarrow \tau" ("\diamond4<sub>0</sub>")
                       axiomatization where
                       ax IO1 : "monotone ♦¹₀" and
                       ax IO2 : "regular dia ♦²₀" and
                       ax IO3i : "monotone ♦³₀" and
                       ax I03ii : "(\diamondsuit^3, \varphi) \leq \diamondsuit^3, (\varphi \land \diamondsuit^3, \varphi)" and
                       ax IO4i : "regular dia ♦⁴₀" and
                       ax I04ii : "(\diamondsuit^4_{\circ} \varphi) \leq \diamondsuit^4_{\circ} (\varphi \wedge \diamondsuit^4_{\circ} \varphi)"
(*Soundness out1*)
lemma IO1top: "\Diamond_0 T \leq T " by (meson settrue def)
lemma IO1SI: "((\diamondsuit^1_0 \ \alpha \le \varphi)) \land (\beta \le \alpha)) \longrightarrow (\diamondsuit^1_0 \beta \le \varphi) " using monotone def ax IO1 by force
lemma IO1W0: "((\diamondsuit^1_0 \ \alpha \le \varphi) \land (\varphi \le \psi)) \longrightarrow (\diamondsuit^1_0 \ \alpha \le \psi) " by simp
lemma IO1AND: "((\diamondsuit^1, \alpha \le \varphi)) \land (\diamondsuit^1, \alpha \le \psi)) \longrightarrow (\diamondsuit^1, \alpha \le (\varphi \land \psi)) " by (simp add: setand def)
(*Soundness out2*)
lemma IO2top: "\diamond^2 T \leq T " by (simp add: settrue def)
lemma IO2SI: "((\diamondsuit^2, \alpha \le \varphi)) \land (\beta \le \alpha)) \longrightarrow (\diamondsuit^2, \beta \le \varphi)" using regular tomonotone
   using monotone def ax IO2 by fastforce
lemma IO2WO: "((\diamondsuit^2, \alpha \leq \varphi)) \land (\varphi \leq \psi)) \longrightarrow (\diamondsuit^2, \alpha \leq \psi)" by simp
lemma IO2AND: "((\diamondsuit^2, \alpha \le \varphi) \land (\diamondsuit^2, \alpha \le \psi)) \longrightarrow (\diamondsuit^2, \alpha \le (\varphi \land \psi)) " by (simp add: setand def)
lemma IO2OR: "((\diamondsuit^2, \alpha \leq \varphi) \land (\diamondsuit^2, \beta \leq \varphi)) \longrightarrow (\diamondsuit^2, (\alpha \lor \beta) \leq \varphi)"
   using ax IO2 regular dia def regular dia def setor def by auto
(*Soundness out3*)
lemma IO3top: "\diamondsuit^3, T \leq T " by (simp add: settrue def)
lemma IO3SI: "((\diamondsuit^3, \alpha \le \varphi) \land (\beta \le \alpha)) \longrightarrow (\diamondsuit^3, \beta \le \varphi)"
   using monotone def ax IO3i by fastforce
lemma IO3WO: "((\diamondsuit^3, \alpha \le \varphi) \land (\varphi \le \psi)) \longrightarrow (\diamondsuit^3, \alpha \le \psi) " by simp
lemma IO3AND: "((\diamondsuit^3, \alpha \le \varphi) \land (\diamondsuit^3, \alpha \le \psi)) \longrightarrow (\diamondsuit^3, \alpha \le (\varphi \land \psi))" by (simp add: setand def)
lemma IO3CT: "((\diamondsuit^3, \alpha < \varphi) \land (\diamondsuit^3, (\alpha \land \varphi) < \psi)) \longrightarrow (\diamondsuit^3, \alpha < \psi)"
   using monotone def ax IO3i ax IO3ii setand def by (smt (verit, ccfv threshold))
```

consts I01 :: " $\tau \Rightarrow \tau$ " (" $\diamondsuit$ 1<sub>0</sub>")



## 



#### Photo from silicon Luxembourg

#### **Student Robots Clash: Cooperation Crisis in Schools**

In 2035, schools face new challenges as student-owned robots struggle to cooperate, leading to conflicts and disruptions. Educators are working to address the growing issue of robot rivalry in classrooms.



### LogikEy Team

\*Obligation



\*Permission



\*Prohibition



\*Design Platform V



\*Online Update

#### Generative Agents: Interactive Simulacra of Human Behavior

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## NEURIPS 2024: THE CONCORDIA CONTEST



ORGANIZED BY: Concordia Contest Team
CURRENT PHASE ENDS: November 1, 2024 At 1:00 AM GMTCURRENT SERVER TIME: October 10, 2024 At 11:28 AM GI

Docker image: ihsaanullah/concordia:latest

Competition Report: https://www.cooperativeai.com/contests/concordia-2024

Oct 2024 Nov 2024

RL 2020: Humans punish and reward Al agents.

By 2030: Al agents will be capable of punishing each other.

Deception: Al agents deceive other Al agents and humans to avoid punishment or receive more rewards.

## LogikEy for Organizations:

Please contact <u>ali.farjami@uni.lu</u> for dealing prices and more discussions!

Special Offer for Kindergartens